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RE-PROXIMITIES AS FIXED POINTS OF AN OPERATOR ON PSEUDO PROXIMITIES

Monique CHICOURRAT

Abstract. We show that the RE-proximities of [7] can be obtained as images of some natural operator defined on pseudo proximities. Considering the one-to-one correspondence between closed graph relations on the ultrafilter space (the so-called "nasses" of [10]) and pseudo proximities, this operator is indeed the correspondent of an idempotent operator defined on nasses and based upon the concepts of equivalence kernel and domain of a nasse.

0 - Introduction.

The results of this paper were presented in [2] and announced in [3] without proofs. Some proofs are different from those in [2].

The aim of this paper is to show that RE-proximities can be constructed as fixed points of some natural operator on the set of *pseudo proximities*.

We denote by $\mathcal{P}(X)$ the power set of a set X.

Definition. Given any relation δ on $\mathcal{P}(X)$, we call a *Riesz extension of* δ any topological space Y containing a copy of X and satisfying the two following conditions:

- (E1) $A\delta B \iff \overline{A} \cap \overline{B} \neq \emptyset \text{ for all } A, B \in \mathcal{P}(X)$
- (E2) $\{x\}\delta A \iff x \in \overline{A} \quad \text{for all} \quad x \in X, \ A \in \mathcal{P}(X).$

Of course, any relation δ admitting a Riesz extension is a *proximity on* X, i.e. δ is a symmetric relation on $\mathcal{P}(X)$ satisfying (P1)-(P3).

- (P1) \emptyset (non- δ) A
- (P2) $A\delta(B \cup C) \iff A\delta B$ or $A\delta C$
- $(P3) A \cap B \neq \emptyset \implies A\delta B$

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Furthermore if δ admits a T_1 -extension, then δ is separated i.e.

$$(P4) \{x\} \delta \{y\} \implies x = y$$

Proximities admitting T_1 -extensions are known. First, it is obvious that such a proximity is a separated LO-proximity, i.e. satisfying the following condition:

(LO)
$$\operatorname{adh}_{\delta}(A) \ \delta \ \operatorname{adh}_{\delta}(B) \implies A\delta B$$

where $\operatorname{adh}_{\delta}(A) = \{x \in X \ / \ \{x\}\delta A\}.$

Conversely, Gagrat, Naimpally [8] and finally Thron [15] proved that any separated LO-proximity admits a compact T_1 -extension to the set Γ_{δ} of maximal " δ -bunches" with the "absorption topology" : $\varphi : x \in X \mapsto \delta(\{x\}) \in \Gamma_{\delta}$ is an embedding of X into Γ_{δ} .

In fact, for any LO-proximity δ on X, define $\hat{\Gamma} = X \cup (\Gamma_{\delta} \setminus \Gamma_{X})$ where $\Gamma_{X} = \{\delta(\{x\}) : x \in X\}$. Consider the following modification of the absorption topology (similar to constructions of Császár): the family $\{\hat{\Gamma}(A) : A \subset X\}$ where

$$\hat{\Gamma}(A) = \{x \in X : A \in \delta(\{x\})\} \cup \{\mathcal{G} \in (\Gamma_{\delta} \backslash \Gamma_X) : A \in \mathcal{G}\}$$

is a closed basis for a topology on $\hat{\Gamma}$ and this topological space provides a compact Riesz extension for the LO-proximity δ .

This proves that proximities admitting Riesz extensions are exactly LO-proximities and that they actually admit a compact extension.

On the other hand, Smirnov [14] proved in 1952 that Efremovič proximities are those proximities admitting T_2 -compact extensions. These are separated EF-proximities (in the sense of [8]), i.e. satisfying the following

(EF)
$$A (non-\delta) B \iff \exists C \subset X : A (non-\delta) C \text{ and } (X \setminus C) (non-\delta) B$$

Constructions given by Smirnov and then by Leader are also absorption topologies on some sets of filters (ends [14]) or grills (clusters [11]).

In the following sections we will be interested in the RE-proximities:

Definition. [7] An RE-proximity is a proximity admitting some regular Riesz extension.

These proximities were characterized by Császár in 1986; in particular, he constructed regular Riesz extensions to the set of some filters (round and compressed filters).

In 1990, Fougères (*) described a new method (fairly similar to Samuel's method [13]) to construct compact extensions for Efremovič proximities, using ultrafilters and based upon the so-called *nasses* of [10].

^(*) Personal communication

This method is of real interest in the study of RE-proximities [5] (also in the study of "RI-proximities" [4]). It naturally yields new characterizations for these proximities and also a generalization of known results for ends or clusters and EF-proximities [12] to round and compressed filters and RE-proximities.

Moreover, we will prove here that this method also shows clearly that RE-proximities are fixed points (or images) of some natural idempotent operator on pseudo proximities.

In the first section, we will recall some preliminary facts and in particular the one-toone correspondence between nasses and pseudo proximities.

Then (section 2), using the concept of equivalence kernel of a nasse, we naturally introduce and study some operator \mathcal{E} on the set of nasses and also the correspondent operator \mathcal{R} defined on the set of pseudo proximities.

Next, we will prove in section 3 that RE-proximities are exactly those proximities which are fixed points of the operator \mathcal{R} .

Finally, in section 4, we link round and compressed filters with ultrafilters. Then one gets a direct construction of the operator \mathcal{R} on pseudo proximities in terms of round and compressed filters.

1 - Preliminaries.

A filter on a set X is a non empty subset \mathcal{F} of $\mathcal{P}(X)$ such that

$$A \cap B \in \mathcal{F} \iff A \in \mathcal{F} \text{ and } B \in \mathcal{F}$$

and a grill on X is a subset \mathcal{G} of $\mathcal{P}(X)$ such that

$$\emptyset \notin \mathcal{G}$$
 and $(A \cup B \in \mathcal{G} \iff A \in \mathcal{G} \text{ or } B \in \mathcal{G}).$

A proper filter will be any filter different from $\mathcal{P}(X)$.

Recall first that the map $g: A \subset \mathcal{P}(X) \mapsto A^{\#} \subset \mathcal{P}(X)$ where

$$\mathcal{A}^{\#} = \{ B \subset X : \forall A \in \mathcal{A}, \ A \cap B \neq \emptyset \}$$

is one-to-one between filters and grills on X, the inverse map is again g ([6], [16]).

We denote by $\Omega(X)$ the set of ultrafilters on X. Then recall also that for any filter \mathcal{F} and grill \mathcal{G} on X,

$$\mathcal{F} \subset \mathcal{G} \iff \exists \mathcal{U} \in \Omega(X) : \mathcal{F} \subset \mathcal{U} \subset \mathcal{G}$$

and for $A \in \mathcal{P}(X)$,

$$A \in \mathcal{G} \iff \exists \mathcal{U} \in \Omega(X) : A \in \mathcal{U} \subset \mathcal{G}.$$

We consider the usual separated compact topology on $\Omega(X)$. The family $\{\Omega(A) : A \subset X\}$ where $\Omega(A)$ is the set of ultrafilters on X containing A is an open and closed basis for this topology, it is closed under finite intersection and finite union.

Thus, if \mathbb{F}_1 and \mathbb{F}_2 are two disjoint closed sets of $\Omega(X)$ then there exists $A \subset X$ such that $\mathbb{F}_1 \subset \Omega(A)$ and $\mathbb{F}_2 \subset \Omega(X) \setminus \Omega(A) = \Omega(X \setminus A)$.

For $\mathbb{H} \subset \Omega(X)$, $\overline{\mathbb{H}}$ denotes the closure of \mathbb{H} in $\Omega(X)$, and if $\Theta \subset \Omega(X)^2$, we write $\overline{\Theta}$ for the closure of Θ for the product topology on $\Omega(X)^2$.

It is known that there is a bijection between filters (or grills) on X and closed subsets of $\Omega(X)$. Any filter \mathcal{F} (or any grill $\mathcal{F}^{\#}$) on X and the closed subset

$$\mathbb{F} = \{ \mathcal{U} \in \Omega(X) : \mathcal{F} \subset \mathcal{U} \} = \{ \mathcal{U} \in \Omega(X) : \mathcal{U} \subset \mathcal{F}^{\#} \}$$

will be said to be associated. And we have

$$\mathcal{F} = \cap \{\mathcal{U} : \mathcal{U} \in \mathbb{F}\} \quad \text{ and } \quad \mathcal{F}^{\#} = \cup \{\mathcal{U} : \mathcal{U} \in \mathbb{F}\}.$$

Furthermore, for any $\mathbb{H} \subset \Omega(X)$, we have $\overline{\mathbb{H}} = \{ \mathcal{U} \in \Omega(X) : \mathcal{F} \subset \mathcal{U} \}$ where $\mathcal{F} = \cap \{ \mathcal{U} : \mathcal{U} \in \mathbb{H} \}$.

¿From now on, we identify any point $x \in X$ with the ultrafilter centered at x. Thus we can write $X \subset \Omega(X)$ and for $A \subset X$, the closure \overline{A} of A in $\Omega(X)$ is $\Omega(A)$. In particular, X is a dense subset of $\Omega(X)$.

Definition 1.1.

- 1. Following Haddad's terminology, we call nasse on X any closed graph relation on $\Omega(X)$.
- 2. Any relation δ on $\mathcal{P}(X)$ such that δ and δ^- satisfy (P1) and (P2) will be called pseudo proximity on X.

In the following theorem, we summarize preliminary results. These were given by Fougères and they are a small generalization of results due to Haddad [10]. One can consult [1] for basic facts on Galois connections. Proofs can be found in [4] and [5].

Theorem 1.1. For any relation δ on $\mathcal{P}(X)$, consider

$$\varphi(\delta) = \{ (\mathcal{U}, \mathcal{V}) \in \Omega(X)^2 : \mathcal{U} \times \mathcal{V} \subset \delta \}$$

and for any relation Θ on $\Omega(X)$, define

$$\psi(\Theta) = \bigcup \{ \mathcal{U} \times \mathcal{V} : (\mathcal{U}, \mathcal{V}) \in \Theta \}$$

$$i.e. \quad A \ \psi(\Theta)B \quad \textit{iff} \quad \Theta(\overline{A}) \cap \overline{B} \ \neq \emptyset \ (\textit{or} \ (\overline{A} \times \overline{B}) \cap \Theta \neq \emptyset).$$

1. (φ, ψ) is a Galois connection i.e. φ and ψ are monotone and

$$\Theta \subset \varphi(\delta) \iff \psi(\Theta) \subset \delta.$$

Thus the restriction $\varphi: \operatorname{Im}\psi \to \operatorname{Im}\varphi$ is the inverse of $\psi: \operatorname{Im}\varphi \to \operatorname{Im}\psi$.

- 2. For a relation Θ on $\Omega(X)$, we have $\varphi(\psi(\Theta)) = \overline{\Theta}$.
- 3. $\Theta \in \operatorname{Im} \varphi$ iff Θ is a nasse on X. $\delta \in \operatorname{Im} \psi$ iff δ is a pseudo proximity on X.
- 4. For any pseudo proximity δ on X,

$$\delta$$
 is symmetric iff $\varphi(\delta)$ is symmetric, δ satisfies (P3) iff $\varphi(\delta)$ is reflexive.

It follows in particular from this theorem that there is a one-to-one correspondence between proximities and reflexive and symmetric closed relations on $\Omega(X)$.

Definition 1.2. For any pseudo proximity δ , the relation

$$\Delta = \{ (\mathcal{U}, \mathcal{V}) \in \Omega(X)^2 : \mathcal{U} \times \mathcal{V} \subset \delta \}$$

is said to be the nasse of δ .

Remark 1.1. The nasse of δ is a closed subset Δ of $\Omega(X)^2$ and for any closed subset \mathbb{F} of $\Omega(X)$, $\Delta(\mathbb{F})$ is closed in $\Omega(X)$.

Remark 1.2. For any pseudo proximity δ with associated nasse Δ , we have

$$A\delta B$$
 iff $\Delta(\overline{A}) \cap \overline{B} \neq \emptyset$.

Finally, we recall the following nice characterization given by Haddad ([9], [10]) for EF-proximities. This characterization is the main fact leading to a simple compact extension for EF-proximities to the set of ultrafilters [5].

Theorem 1.2. A (pseudo) proximity δ with nasse Δ satisfies condition (EF) iff Δ is an equivalence relation on $\Omega(X)$.

2 - The operators \mathcal{E} and \mathcal{R}

The following lemmas introduce the concepts of equivalence kernel and domain of a symmetric relation. (The proofs are left to the reader).

Firstly, we will show that these notions generate idempotent operators on symmetric nasses and on symmetric pseudo proximities.

Lemma 2.1. Let R be a symmetric relation on a set E and define

$$\mathbb{D}(R) = \{x \in E : R(x) \neq \emptyset \text{ and for any } y \in R(x), R(y) = R(x)\}.$$

Then $R(\mathbb{D}(R)) = \mathbb{D}(R)$ and R is an equivalence relation on $\mathbb{D}(R)$. Furthermore, $\mathbb{D}(R)$ is the largest subset of E with these two properties.

Definition 2.1. (With the above notations).

The set $\mathbb{D}(R)$ is called the equivalence domain of R and $R' = R \cap \mathbb{D}(R)^2$ is its equivalence kernel. An equivalence class of R' will be said equivalence class of R.

Lemma 2.2. For any symmetric relation R on E, a non void subset A of E is an equivalence class of R iff $A \times A \subset R$ and R(A) = A (or $R(A) \subset A$).

Lemma 2.3. If R is a symmetric relation on E and $F \subset E$, consider the relation $R_F = R \cap F^2$ on F. Then $F \cap \mathbb{D}(R) \subset \mathbb{D}(R_F)$ and $R' \cap F^2 \subset R'_F$.

Remark 2.1. For any $x \in E$, we have :

$$x \in \mathbb{D}(R) \iff (R(x) \neq \emptyset, R(x) \times R(x) \subset R, \text{ and } R(R(x)) = R(x)).$$

Remark 2.2. The equivalence kernel R' of R is not necessarily a maximal equivalence relation contained in R. It can even be empty. For example, if $E = \{1, 2, 3\}$, and

$$x R y \iff (x = y) \text{ or } (x = 1 \text{ and } y \in \{2,3\}) \text{ or } (y = 1 \text{ and } x \in \{2,3\}).$$
 Then $R' = \emptyset$.

Definition 2.2.

- 1. For any symmetric nasse Δ , let Δ' be the equivalence kernel of Δ and define $\mathcal{E}(\Delta)$ to be the (symmetric) nasse $\overline{\Delta'}$.
- 2. For any symmetric pseudo proximity δ with symmetric nasse Δ , we define $\mathcal{R}(\delta)$ to be the symmetric pseudo proximity on X whose nasse is $\mathcal{E}(\Delta) = \overline{\Delta'}$.

Proposition 2.4.

- 1. For any symmetric nasse Δ , $\mathcal{E}(\Delta)$ is also symmetric and : $\mathcal{E}(\Delta) \subset \Delta$ and $\mathcal{E}(\mathcal{E}(\Delta)) = \mathcal{E}(\Delta)$.
 - 2. For any symmetric pseudo proximity δ : $\mathcal{R}(\delta) \subset \delta$ and $\mathcal{R}(\mathcal{R}(\delta)) = \mathcal{R}(\delta)$.

Proof. First observe that if Δ' is the equivalence kernel of Δ then $\mathcal{E}(\Delta) = \overline{\Delta'}$ is symmetric since Δ' is.

Of course, $\mathcal{E}(\Delta) \subset \Delta$ since $\Delta' \subset \Delta$. Thus, to prove that \mathcal{E} is idempotent we show that $\mathcal{E}(\Delta) \subset \mathcal{E}(\mathcal{E}(\Delta))$.

Actually, we have $\mathbb{D}(\Delta) \subset \mathbb{D}(\mathcal{E}(\Delta))$ since for $\mathcal{U} \in \mathbb{D}(\Delta)$, $\Delta(\mathcal{U}) = \Delta'(\mathcal{U}) = \overline{\Delta'}(\mathcal{U})$ and then

$$\Delta(\mathcal{U}) \times \Delta(\mathcal{U}) \subset \Delta' \subset \overline{\Delta'} \quad \text{ and } \quad \overline{\Delta'} \left(\overline{\Delta'} \left(\mathcal{U} \right) \right) \subset \Delta(\Delta(\mathcal{U})) \subset \Delta(\mathcal{U}) = \overline{\Delta'} \left(\mathcal{U} \right).$$

Hence, $\Delta' = \Delta' \cap \mathbb{D}(\Delta)^2 \subset \overline{\Delta'} \cap \mathbb{D}(\Delta)^2 \subset \overline{\Delta'} \cap \mathbb{D}(\mathcal{E}(\Delta))^2 = (\mathcal{E}(\Delta))'$ and thus $\mathcal{E}(\Delta) \subset \mathcal{E}(\mathcal{E}(\Delta))$.

Finally, assertion 2. of the proposition is true since the operators φ and ψ of theorem 1.1 are monotone.

Remark 2.3. If $\mathcal{E}(\Delta)$ is reflexive so is Δ and thus if $\mathcal{R}(\delta)$ is a proximity, so is δ . The converse is not true. Furthermore, the operators \mathcal{E} and \mathcal{R} are neither increasing nor decreasing operators.

Actually, assume that X contains at least three points $x, y, z \in X$. Let Δ_1 be the diagonal of $\Omega(X)^2$ and consider $\Delta_2 = \Delta_1 \cup \{(x,y),(y,x),(x,z),(z,x)\}$ and $\Delta_3 = \Delta_2 \cup \{(y,z),(z,y)\}$. These relations are symmetric nasses on X, we have $\mathcal{E}(\Delta_1) = \Delta_1$ and $\mathcal{E}(\Delta_3) = \Delta_3$. And since any singleton of X^2 is open in $\Omega(X)^2$, we have $\mathcal{E}(\Delta_2) = \Delta_2' = \{(\mathcal{U},\mathcal{U}) : \mathcal{U} \in \Omega(X) \setminus \{x,y,z\}\}$.

Now observe that Δ_2 is reflexive, but $\mathcal{E}(\Delta_2)$ is not. And we have $\Delta_1 \subset \Delta_2 \subset \Delta_3$ but $\mathcal{E}(\Delta_1) \not\subset \mathcal{E}(\Delta_2)$ and $\mathcal{E}(\Delta_3) \not\subset \mathcal{E}(\Delta_2)$.

Proposition 2.5. For any proximity δ with nasse Δ , the pseudo proximity $\mathcal{R}(\delta)$ is a proximity iff $X \subset \mathbb{D}(\Delta)$.

Proof. We prove that for any symmetric nasse Δ , $\mathcal{E}(\Delta)$ is reflexive iff $X \subset \mathbb{D}(\Delta)$. Actually, for any $x \in X$, since $\{(x, x)\}$ is open in $\Omega(X)^2$, we have the following:

$$x \in \mathbb{D}(\Delta) \iff (x,x) \in \Delta' \iff (x,x) \in \overline{\Delta'}.$$

Furthermore, $\{(x,x):x\in X\}$ is dense in the diagonal of $\Omega(X)^2$. Thus, $X\subset \mathbb{D}(\Delta)$ iff $\mathcal{E}(\Delta)=\overline{\Delta'}$ is reflexive.

3. RE-PROXIMITIES AS IMAGES OF THE OPERATOR \mathcal{R}

In this section, we use results and constructions presented also in [5].

The aim of this section is to prove that RE-proximities are exactly the proximities which are the fixed points (or images) of \mathcal{R} :

Theorem 3.1. If δ is a proximity on X with associated nasse Δ , let Δ' be the equivalence kernel of Δ and $\mathbb{D}(\Delta)$ be the equivalence domain of Δ .

Then, δ is an RE-proximity iff one of the following conditions is satisfied. (By theorem 1.1., these conditions are equivalent.)

(RE1)
$$\mathcal{E}(\Delta) = \Delta$$

(RE2)
$$\mathcal{R}(\delta) = \delta$$
, or equivalently : $(A\delta B \iff \Delta'(\overline{A}) \cap \overline{B} \neq \emptyset)$.

Furthermore, if these conditions are satisfied, then δ admits a regular Riesz extension to $\mathbb{D}(\Delta)$.

Theorem 3.2. Let δ be any proximity with nasse Δ .

1. Considering Δ' as a relation on $\Omega(X)$, define the following operator on $\mathbb{D}(\Delta)$: $c: \mathbb{H} \subset \mathbb{D}(\Delta) \mapsto \Delta'(\overline{\mathbb{H}}) \subset \mathbb{D}(\Delta).$

This operator is a regular topological closure on $\mathbb{D}(\Delta)$.

2. If $\mathcal{R}(\delta)$ is a proximity then $\mathcal{R}(\delta)$ is an RE-proximity and $(\mathbb{D}(\Delta), c)$ is a regular extension of $\mathcal{R}(\delta)$.

Lemma 3.3. Suppose that E is a topological space and R is a symmetric relation on E such that for all closed subsets F of E, R(F) is closed in E. Let R' be the equivalence kernel of R.

Then the operator $c: H \subset \mathbb{D}(R) \mapsto R'(\overline{H}) \subset \mathbb{D}(R)$ defines a topology on $\mathbb{D}(R)$ and if E is a separated normal space, then this topology is regular.

Proof. Remark that for $H \subset E$, we have $R'(H) = R'(H \cap \mathbb{D}(R)) = R(H) \cap \mathbb{D}(R)$ and if $H \subset \mathbb{D}(R)$, R'(H) = R(H).

That c is a topological closure on $\mathbb{D}(R)$ is now based on the following observation: for $H \subset \mathbb{D}(R)$, we have $R'(\overline{H}) = R'(\overline{H} \cap \mathbb{D}(R)) = R(\overline{H}) \cap \mathbb{D}(R)$, then c(H) is finally the image of the closure of H in the subspace $\mathbb{D}(R)$ by the equivalence relation R' and it is closed in $\mathbb{D}(R)$.

To prove the regularity of $(\mathbb{D}(R), c)$, take $x \in \mathbb{D}(R)$, $H \subset \mathbb{D}(R)$ such that $x \notin R'(\overline{H})$. Then, $R'(x) \cap R'(\overline{H}) = \emptyset$, and hence, R(x) and $R(\overline{H})$ are disjoint closed subsets of E (recall that $R'(x) = R(x) \subset \mathbb{D}(R)$).

Take closed sets F_1 and F_2 in E such that $F_1 \cup F_2 = E$ and $R(x) \cap F_1 = \emptyset$, $R(\overline{H}) \cap F_2 = \emptyset$. Then $R'(F_1) \cup R'(F_2) = \mathbb{D}(R)$, $x \notin R'(F_1)$, and $R'(\overline{H}) \cap R'(F_2) = \emptyset$. Remark now that $R'(F_i) = R'(\overline{F_i \cap \mathbb{D}(R)})$ and the proof is complete.

Proof of theorem 3.2. ¿From lemma 3.3, the operator c defines a regular topology on $\mathbb{D}(\Delta)$. To prove the assertion 2., first, let us notice that Δ' is an equivalence relation on $\mathbb{D}(\Delta)$ and thus by theorem 1.1, we have

$$A \mathcal{R}(\delta) B \iff (\overline{A} \times \overline{B}) \cap \Delta' \neq \emptyset \iff \Delta'(\overline{A}) \cap \Delta'(\overline{B}) \neq \emptyset.$$

Hence, if $\mathcal{R}(\delta)$ is a proximity, then $X \subset \mathbb{D}(\Delta)$ (by proposition 2.5.), and we have

(E1) :
$$A \mathcal{R}(\delta) B \iff c(A) \cap c(B) \neq \emptyset$$

(E2) :
$$\{x\} \mathcal{R}(\delta) A \iff (\{x\} \times \overline{A}) \cap \Delta' \neq \emptyset \iff x \in \Delta'(\overline{A}) \iff x \in c(A)$$

and $(\mathbb{D}(\Delta), c)$ is a regular Riesz extension of $\mathcal{R}(\delta)$.

Theorem 3.2. proves that if the conditions (RE) are satisfied then δ is an RE-proximity. The two next lemmas prove that the converse is true.

Lemma 3.4. If Y is a topological space, let π be the proximity defined on Y by $(A\pi B \iff c(A) \cap c(B) \neq \emptyset)$ where c(A) is the closure of A in Y. Denote by Δ the nasse of π . If Y is regular, then $Y \subset \mathbb{D}(\Delta)$ and π satisfies (RE2) (i.e. $\mathcal{R}(\pi) = \pi$).

Proof. To start with, remark that if Y is regular, then $y \in c(A)$ iff $\{y\}\pi A$.

Fix $y \in Y$. To prove that $y \in \mathbb{D}(\Delta)$, we use remark 2.1.

If $(\mathcal{U}, \mathcal{V}) \in \Delta(y) \times \Delta(y)$, for any $(A, B) \in \mathcal{U} \times \mathcal{V}$, we have $\{y\}\pi A$ and $\{y\}\pi B$, and then $A\pi B$. So $(\mathcal{U}, \mathcal{V}) \in \Delta$ and $\Delta(y) \times \Delta(y) \subset \Delta$.

Furthermore, fix $\mathcal{U} \in \Delta(y)$ and $\mathcal{V} \in \Delta(\mathcal{U})$. For any $A \in \mathcal{U}$, we have $y \in c(A)$ and for any $B \in \mathcal{V}$, $c(A) \cap c(B) \neq \emptyset$. Now set $B \in \mathcal{V}$, if $y \notin c(B)$, take an open set O such that $y \in O \subset c(O) \subset Y \setminus c(B)$. Then $y \notin Y \setminus O$, so $Y \setminus O \notin \mathcal{U}$, and also $c(B) \cap c(O) = \emptyset$, and thus $O \notin \mathcal{U}$. This contradicts that \mathcal{U} is a ultrafilter. Therefore, $\mathcal{V} \in \Delta(y)$ and finally, $\Delta(\Delta(y)) = \Delta(y)$.

Hence $Y \subset \mathbb{D}(\Delta)$ and :

$$A\pi B \iff \exists y \in Y : \{y\}\pi A \text{ and } \{y\}\pi B$$

$$\iff \exists y \in Y : \overline{A} \cap \Delta(y) \neq \emptyset \text{ and } \overline{B} \cap \Delta(y) \neq \emptyset$$

$$\iff \Delta'(\overline{A}) \cap \overline{B} \neq \emptyset \quad \text{(because } \Delta(y) \times \Delta(y) \subset \Delta'\text{)}.$$

$$\iff A \mathcal{R}(\pi) B$$

Finally, since $\mathcal{R}(\pi) \subset \pi$, we have $\mathcal{R}(\pi) = \pi$, and condition (RE2) is satisfied.

Lemma 3.5. If π is a proximity on a set Y and $X \subset Y$, let δ be the proximity induced by π on X, i.e.

$$\delta = \pi \cap \mathcal{P}(X)^2.$$

If $\mathcal{R}(\pi) = \pi$ then $\mathcal{R}(\delta) = \delta$.

Proof. Let Δ_Y be the nasse of π , and Δ_X that of δ . We have to prove that if $\mathcal{E}(\Delta_Y) = \Delta_Y$ then $\mathcal{E}(\Delta_X) = \Delta_X$.

Since $X \subset Y$, $\Omega(X)$ can be identified with the open-closed subset $\{\mathcal{U} \in \Omega(Y) : X \in \mathcal{U}\}$ of $\Omega(Y)$. Now one can easily verify that $\Delta_X = \Delta_Y \cap \Omega(X)^2$.

So, by Lemma 2.3, $\Delta'_Y \cap \Omega(X)^2 \subset \Delta'_X$. But $\Omega(X)^2$ is open in $\Omega(Y)^2$; for this reason, the density of Δ'_Y in Δ_Y implies that of $\Delta'_Y \cap \Omega(X)^2$ in $\Delta_Y \cap \Omega(X)^2$, and thus, that of Δ'_X in Δ_X .

Remark 3.1. It can be shown that any separated RE-proximity actually admits a separated regular Riesz extension [5]: if δ satisfies one of the "(RE)-conditions" and if it is separated, then the quotient space $\mathbb{D}(\Delta)/\Delta'$ provides a separated and regular Riesz extension of δ .

Remark 3.2. We finally remark that of course the constructions and results of Theorem 3.2 and of the preceding remark are available for EF-proximities and in fact, we obtain as a corollary known results for these proximities [5]: if δ is an EF-proximity with nasse Δ , then $\mathbb{D}(\Delta) = \Omega(X)$ and $\Delta' = \Delta$; the operator $c : \mathbb{H} \subset \Omega(X) \mapsto \Delta(\overline{\mathbb{H}}) \subset \Omega(X)$ is a compact topological closure on $\Omega(X)$ and $\Omega(X)$, c is a Riesz extension of δ ; furthermore if δ is separated, then the separated quotient space $\Omega(X)/\Delta$ provides a separated compact Riesz extension of δ .

4. The operator \mathcal{R} and round and compressed filters

In this section we will prove that the pseudo proximity $\mathcal{R}(\delta)$ can be directly defined from δ , using the link between round and compressed filters and ultrafilters.

We need some preliminary notations.

Let δ be any pseudo proximity on X and Δ its associated nasse.

Notations. Let γ be the relation defined on $\mathcal{P}(X)$ by

$$A\gamma B \iff A (non-\delta) (X\backslash B).$$

Again we denote by $\gamma(A)$ the set $\{B \subset X : A\gamma B\}$.

If $\mathcal{A} \subset \mathcal{P}(X)$, we define

$$\gamma(\mathcal{A}) = \bigcup \{\gamma(A) : A \in \mathcal{A}\} \quad \text{and} \quad \mathcal{A}^{\delta} = \bigcap \{\delta(A) : A \in \mathcal{A}\}.$$

Remark that for any filter \mathcal{F} on X, $\gamma(\mathcal{F})$ is a filter and \mathcal{F}^{δ} is a grill on X.

The following lemma is a small generalization of a result given in [10] for proximities. The proof can be found also in [5].

Lemma 4.1. For any filter \mathcal{F} on X with associated closed $\mathbb{F} \subset \Omega(X)$, $\gamma(\mathcal{F})^{\#} = \mathcal{F}^{\delta}$ and the closed subset of $\Omega(X)$ associated with the filter $\gamma(\mathcal{F})$ (or with the grill \mathcal{F}^{δ}) is $\Delta(\mathbb{F})$.

Recall now some definitions. The definition of δ -cluster is due to Leader [11].

Definition 4.1. ([7], [16]) A proper filter \mathcal{F} on X is said to be δ -round iff $\gamma(\mathcal{F}) = \mathcal{F}$. We say that \mathcal{F} is δ -compressed or $\mathcal{F}^{\#}$ is a δ -clan iff $\mathcal{F}^{\#} \times \mathcal{F}^{\#} \subset \delta$. A grill \mathcal{G} on X is called a δ -cluster iff $\mathcal{G} = \mathcal{G}^{\delta}$.

It is easy to see that any δ -cluster is a maximal δ -clan (the converse is false [2]) and note that if δ is a proximity, then for any $\mathcal{U} \in \Omega(X)$, $\mathcal{U} \subset \mathcal{U}^{\delta}$ and thus $(\mathcal{U}^{\delta})^{\delta} \subset \mathcal{U}^{\delta}$. It follows that the grill \mathcal{U}^{δ} is a δ -clan iff \mathcal{U}^{δ} is a δ -cluster.

Recall that for any filter \mathcal{F} , with associated closed $\mathbb{F} \subset \Omega(X)$, $\mathcal{F}^{\#}$ is the union of the ultrafilters of \mathbb{F} . So, using Lemma 4.1., one gets the following result.

Proposition 4.2. For any proper filter \mathcal{F} on X with associated closed $\mathbb{F} \subset \Omega(X)$: \mathcal{F} is δ -round iff $\mathbb{F} = \Delta(\mathbb{F})$; and \mathcal{F} is δ -compressed iff $\mathbb{F} \times \mathbb{F} \subset \Delta$.

In particular, the associated closed $\mathbb{F} \subset \Omega(X)$ of a ultrafilter \mathcal{U} is $\{\mathcal{U}\}$. So by Lemma 4.1, $\Delta(\mathcal{U})$ is the closed subset of $\Omega(X)$ associated to the filter $\gamma(\mathcal{U})$ and therefore with Lemma 2.2, we can deduce the following theorem.

Theorem 4.3. For any symmetric pseudo proximity δ :

- 1. $\mathcal{U} \in \mathbb{D}(\Delta) \iff \gamma(\mathcal{U}) \text{ is } \delta\text{-round compressed.}$
- 2. For any proper filter \mathcal{F} on X, with associated closed $\mathbb{F} \subset \Omega(X)$, \mathcal{F} is δ -round compressed iff there exists $\mathcal{U} \in \mathbb{D}(\Delta)$ such that $\mathbb{F} = \Delta(\mathcal{U})$ and then $\mathcal{F} = \gamma(\mathcal{U})$ and $\mathcal{F}^{\#} = \mathcal{U}^{\delta}$.

Thus for any proximity δ , if \mathcal{F} is a δ -round compressed filter, then $\mathcal{F}^{\#}$ is a δ -cluster.

This result now yields a direct construction of $\mathcal{R}(\delta)$ from δ and hence gives also characterizations of RE-proximities. The characterization (RE4) was given (in terms of the "order" γ) by Császár [7].

For any symmetric pseudo-proximity δ with associated nasse Δ , denote by Δ' the equivalence kernel of Δ . Note that for any $A \in \mathcal{P}(X)$, and any $\mathcal{U} \in \Omega(X)$, we have :

 $\mathcal{U} \in \Delta'(\overline{A})$ iff $\mathcal{U} \in \mathbb{D}(\Delta)$ and $A \in \mathcal{U}^{\delta}$. Since Δ' is an equivalence relation on $\mathbb{D}(\Delta)$, using the characterization of a pseudo proximity with its associated nasse (theorem 1.1), we deduce:

Theorem 4.4.

- 1. $A\mathcal{R}(\delta)B \iff \exists \mathcal{U} \in \mathbb{D}(\Delta) : A, B \in \mathcal{U}^{\delta}$ $\iff there \ exists \ a \ \delta\text{-round compressed filter} \ \mathcal{F} \ such \ that \ A, B \in \mathcal{F}^{\#}.$
- **2.** A proximity δ is an RE-proximity iff one of the following equivalent statements is satisfied:
- (RE3) $A\delta B \iff \exists \mathcal{U} \in \mathbb{D}(\Delta) : A, B \in \mathcal{U}^{\delta}$
- (RE4) $A\delta B \iff there \ exists \ a \ \delta$ -round compressed filter \mathcal{F} such that $A, B \in \mathcal{F}^{\#}$.

Remark 4.1. Note that theorem 4.3 is a generalization of results available for EF-proximities ([12], [16], [2]). As a matter of fact, the following assertions are equivalent if Δ is an equivalence relation and $\mathbb{F} \subset \Omega(X)$:

- 1. \mathbb{F} is maximal for the property $\mathbb{F} \times \mathbb{F} \subset \Delta$
- 2. **F** is minimal for the property $\mathbb{F} = \Delta(\mathbb{F})$
- 3. $\mathbb{F} \times \mathbb{F} \subset \Delta$ and $\mathbb{F} = \Delta(\mathbb{F})$
- 4. There exists $\mathcal{U} \in \Omega(X)$ such that $\mathbb{F} = \Delta(\mathcal{U})$.

Now let δ be an EF-proximity and \mathcal{F} some filter on X. Since minimality for the filter \mathcal{F} (and maximality for the associated grill) means maximality for the associated closed \mathbb{F} , and vice-versa, then the following properties are equivalent [5]:

- 1. $\mathcal{F}^{\#}$ is a maximal δ -clan
- 2. \mathcal{F} is a minimal δ -compressed filter
- 3. \mathcal{F} is a maximal δ -round filter
- 4. \mathcal{F} is a δ -round compressed filter
- 5. $\mathcal{F} = \gamma(\mathcal{U})$ for some $\mathcal{U} \in \Omega(X)$ (and $\mathcal{F}^{\#} = \mathcal{U}^{\delta}$)
- 6. $\mathcal{F}^{\#}$ is a δ -cluster.

Remark 4.2. For any proximity δ , let us denote by Λ_{δ} the set of δ -round and compressed filters. It is shown in [5] that for any separated RE-proximity δ , one gets a (separated) regular Riesz extension of δ to the set Λ_{δ} . This extension is equivalent (in the same way that compactifications of a completely regular space are equivalent) to $\mathbb{D}(\Delta)/\Delta'$, and to the Riesz extension constructed by Császár.

But also in the general case of RE-proximities, it can be proved that the regular Riesz extension constructed by Császár is a quotient of a subspace of $\mathbb{D}(\Delta)$. Actually, consider

 $\Lambda_X = \{\gamma(\{x\}) : x \in X\} \ (\Lambda_X \text{ is the set of } convergent filters) \text{ and } \hat{\Lambda}_{\delta} = X \cup (\Lambda_{\delta} \setminus \Lambda_X);$ the topology generated by the family $\{\hat{\Lambda}_{\delta}(A) : A \subset X\}$ where $\hat{\Lambda}_{\delta}(A) = \{x \in X : A \in \gamma(\{x\})\} \cup \{\mathcal{F} \in (\Lambda_{\delta} \setminus \Lambda_X) : A \in \mathcal{F}\}$ provides a regular Riesz extension for the RE-proximity δ [7]. One can show that this space is also a quotient of a subspace of $\mathbb{D}(\Delta)$ (see [2]), namely, it is a quotient of $\hat{\mathbb{D}} = (\mathbb{D}(\Delta) \setminus \mathbb{E}_X) \cup X$ where $\mathbb{E}_X = \bigcup_{x \in X} \Delta(x)$.

Remark 4.3. Finally, let us remark that Császár's *preregular systems of filters* can be linked with ultrafilters. These are systems Λ of filters on X inducing a proximity δ as follows, for which $\Lambda \subset \Lambda_{\delta}$:

$$A\delta B \iff \exists \mathcal{F} \in \Lambda : A, B \in \mathcal{F}^{\#}$$

Such a proximity δ is an RE-proximity (theorem 4.4.). If Δ is the nasse of this proximity, then preregular systems of filters Λ inducing δ correspond exactly to subsets \mathbb{E} of $\mathbb{D}(\Delta)$ such that $\mathbb{E} = \Delta(\mathbb{E})$ and $\Delta \cap \mathbb{E}^2$ is dense in Δ . The correspondence is clearly the following: $\mathbb{E} = \{\mathcal{U} \in \mathbb{D}(\Delta) : \gamma(\mathcal{U}) \in \Lambda\}$ and $\Lambda = \{\gamma(\mathcal{U}) : \mathcal{U} \in \mathbb{E}\}.$

Then, the results obtained for the Riesz extensions $\mathbb{D}(\Delta)$, $\mathbb{D}(\Delta)/\Delta'$ and Λ_{δ} can be generalized to \mathbb{E} , $\mathbb{E}/(\Delta \cap \mathbb{E}^2)$ and Λ .

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